

Implementing NRR

An Algorithmic Approach

Outline

1. Adjacency Matrix
2. Procedure / Algorithm - Backward Search
3. Procedure / Algorithm - Forward Search
4. Remarks
5. Future Work
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Adjacency Matrix

- An adjacency matrix for a graphical network contains a '1' in position (i,j) if node i is directly linked with node j via an edge (direct connection)
- We want to use the adjacency matrix to study indirect connections.

				j	
i					

Procedure / Algorithm

- If there is a '0' in (i,j) it means there is no direct connection between nodes i and j

i				0	

Procedure / Algorithm 2 - Backward Search

- Step 1: Consider row i , columns 1 through $j - 1$
- Step 2: If there exists a 1 there, go to Step 3, else go to Step 4'
- Step 3: Note down the column number of the '1.' Say it is column k , where $1 \leq k \leq j$.

		k		j	
i		1		0	

Procedure / Algorithm 3

- Step 4: In column k, consider rows i-1 through 1 in descending order (bigger value to smaller value). If there is a 1 encountered there, note its row number. Say it is l.

	k		j		
l			1		
i			1	0	

	k		j		
l			1		
i			1	0	

Procedure / Algorithm 4

- Step 5: From position (l,k) , scan row l , columns k through j . If position (l,j) is a 1, then stop, since there is an indirect path between i and j . The path is $(i,k) \rightarrow (k,l) \rightarrow (l,j)$. It is a four node path, with three edges.
- Step 6: If no 1 is encountered in Step 4, there is no indirect connection yet found.
- Step 7: If more than one '1' is encountered in Step 4 say on rows $l_1, l_2, \dots$ and Step 5 using the l_1 1 was unsuccessful, repeat Step 5 using the second such '1', namely the one in row l_2 column k .

			k		j	
l			1		1	
i			1		0	

Procedure / Algorithm 5 - Forward Search

- Step 4': Consider row i , columns $j+1$, through n . If there exists a `1' therein, go to Step 5'
- Step 5': Note down its column number. Say it is k' where $j+1 \leq k' \leq n$.
- Step 6': In column k' consider rows $i+1$ through n . If a `1' is encountered, note its row number. Say it is l' .
- Step 7': From position (l', k') scan row l' . If column j has a `1' then we are done.

Remarks

1. Primed steps represent a 'Forward' search
2. Unprimed steps represent a 'Backward' search
3. The adjacency matrix A is sufficient to determine graph 'connectivity' which includes indirect connections.
4. This algorithm will determine at least one "3-edged" indirect path between nodes i and j , if it exists.
5. Versions of this algorithm can be created in future to determine at least one "w-edged" indirect path between nodes i and j if it exists.

Future Work

- Bimodal version
- NRR full implementation:
 - Determine every connected pair (i,j) (direct or indirect)
 - Delete all edges
 - Randomly assign routes for indirectly connected pairs

Allied Work

1. Capacity analysis
 - a. Suppose 1s are replaced by link capacities. Can we use this algorithm to compute max throughput between a pair (i,j) ?
 - b. Compare with max-flow/min-cut theorem
2. Performance analysis
 - a. Instantiate a random graph
 - b. Serve it as input to NRR
 - c. See time taken to perform NRR
 - d. Can random matrix theory plus Lovasz's paper on 'Random walks on graphs' be used to get an analytical result on speed and performance?